

AMENDMENTS TO THE SPECIFICATION

Please amend paragraph 0016 on page 7, as follows:

To implement this scaling, a drop policy could be employed in which four consecutive packets from flow 3 and two packets from flow 4 are dropped. However, it is desirable to employ a random or pseudorandom (or simulated random) drop policy to avoid synchronization of multiple sources. A drop schedule is a function of K and the current load L_i . Thus, the reduction factor is given by $R_i = L_i \times (1-K)$. In one embodiment, the drop factor may be implemented as a ~~cycle~~ cyclic buffer such as, drop buffer 118, in which a "1" indicates a packet drop and a 0 indicates the packet is serviced. After the drop or service decision is made, a drop schedule pointer is advanced to the next location in the drop buffer 118. In one embodiment, all drop schedules are based on a number of packets dropped out of eight possible packets. Thus, eight separate drop buffers may be retained in memory 116 and an appropriate one for the selected scaling factor is selected from the set of existing drop buffers. Packets to be dropped are dropped at the input interface 112.

Please amend paragraph 0018 on page 8, as follows:

At decision block 206, a determination is made if the aggregate load of all of the flows ~~and exceeds~~ a predicted steady state threshold. At decision block 208, if the aggregate flow does not exceed the steady state threshold, the determination is made if the processor is over utilized. If the processor is not over utilized, the system advances to the next time slice and no drop policy is employed at functional block 210. In one embodiment, a time slice is 200ms, other embodiment may employ longer or shorter time slices. If the load is greater than the steady state threshold at decision block 206, at decision block 214 a determination is made if the processor is under utilized. If the processor is under utilized the steady state threshold is raised to more efficiently use the processor. In one embodiment, there is a range in which the steady state threshold may be established. In one such embodiment, the rate at which the steady state threshold is reduced responsive to over utilization exceeds the rate at which the steady state threshold is increased responsive to under utilization. In another embodiment, the steady state threshold is fixed and

unchangeable after manufacture. In such an embodiment, blocks 212-216 are effectively omitted.